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MODEL THEORETIC APPROACH TO TOPOLOGICAL FUNCTORS, II.

by

Jiří Rosický

This paper is a sequel of [6]. Most of results here presented will appear in the forthcoming author's paper [7].

Under a concrete category (\mathcal{A}, U) we will mean a category \mathcal{A} equipped with a faithful functor $U: \mathcal{A} \rightarrow \text{Set}$ satisfying the following two conditions:

- (1) If $A \in \mathcal{A}$, X is a set and $f: UA \rightarrow X$ a bijection, then there is $B \in \mathcal{A}$ and an isomorphism $g: A \rightarrow B$ such that $Ug = f$
- (2) If $A \in \mathcal{A}$ and $f: A \rightarrow A$ is an isomorphism such that Uf is the identity, then f is the identity.

Under a functor $F: (\mathcal{A}, U) \rightarrow (\mathcal{B}, V)$ between concrete categories we will mean a functor $F: \mathcal{A} \rightarrow \mathcal{B}$ such that $V.F = U$.

A type is given by a class of function symbols and a class of relation symbols. Their arities are arbitrary cardinals. The infinitary first-order language $L_{\infty, \omega}(\tau)$ of type τ includes a proper class V of variables and besides the usual logical symbols it admits infinitary conjunctions, disjunctions and quantifiers. A class of sentences of $L_{\infty, \omega}(\tau)$ is called a theory of type τ . We denote by $(\mathcal{A}_\tau, U_\tau)$ or (\mathcal{A}_T, U_T) the concrete category of all τ -structures or T -models resp. These categories need not be legitimate, i.e. they need not form a class. A theory having a representative set of n -ary atomic formulae for each cardinal n will be called normal. If T is normal, then (\mathcal{A}_T, U_T) is a legitimate category and even it is strongly fibre-small in the sense of [1].

If (\mathcal{A}, U) is a concrete category and n a cardinal, then U^n will denote the functor $\text{Set}(n, U-)$. Subfunctors of U^n will be called n -ary relation symbols interpretable in (\mathcal{A}, U) and natural transformations

$U^n \rightarrow U$ n -ary function symbols interpretable in (A, U) . It is motivated by the fact that any relation or function symbol of type \uparrow determines a subfunctor of U^n or a natural transformation $U^n \rightarrow U$ resp. Let \mathcal{S}_U be the collection of all relation and function symbols interpretable in (A, U) . We emphasize that \mathcal{S}_U need not be a type because it need not be a class.

Let $\sim \in \mathcal{S}_U$ be a type. There is a functor $G_\sim: (A, U) \rightarrow (A_\sim, U_\sim)$ such that if $A \in \mathcal{A}$, then the \sim -structure $G_\sim(A)$ has the underlying set UA , the n -ary relation on UA corresponding to $R \in \text{Rel}_n(\sim)$ equals to $R(A)$ and the n -ary function $f: (UA)^n \rightarrow UA$ corresponding to $f \in \text{Fnt}_n(\sim)$ is the component f_A of the natural transformation f . Let T_\sim be the theory of type \sim consisting of all sentences which hold in all \sim -structures $G_\sim(A)$ for $A \in \mathcal{A}$. Clearly we get the functor $G_\sim: (A, U) \rightarrow (A_{T_\sim}, U_{T_\sim})$.

We may restrict ourselves in the formation of T_\sim to some specified kind of sentences. This yields a general method of getting suitable completions or hulls of (A, U) . E.g. (with size conditions aside), if \sim consists of all function symbols from \mathcal{S}_U and $T \in T_\sim$ of all atomic sentences, then T is the Linton's equational theory of U and $G: (A, U) \rightarrow (A_T, U_T)$ is the equational completion of (A, U) (see [5]).

If $A \in \mathcal{A}$, then $R_A(X) = \{Uf / f: A \rightarrow X\}$ defines a subfunctor R_A of U^{UA} . Let $\mathcal{T}_U \in \mathcal{S}_U$ be the type consisting of R_A where A carries over mutually non-isomorphic objects $A \in \mathcal{A}$ such that UA is a cardinal. Then $G_{\mathcal{T}_U}$ is a full embedding and it is important that whenever (A, U) is strongly fibre-small, then $T_{\mathcal{T}_U}$ is normal and (A, U) isomorphic to $(A_{T_{\mathcal{T}_U}}, U_{T_{\mathcal{T}_U}})$.

Further, if T consist of all universal Horn sentences without equality (their specification follows) from $T_{\mathcal{T}_U}$, then (A_T, U_T) is the Mac Neille completion of (A, U) (in the sense of Herrlich [3]). It

proves the conjecture from [6].

Theorem: A concrete category (\mathcal{A}, U) is (absolutely) topological iff it is isomorphic to the category of models of a relational normal universal Horn theory T without equality of some type τ .

Relational means that τ contains relation symbols only and universal Horn theory without equality consists of sentences bearing this name, i.e. arising from formulas $\bigwedge_{i \in I} R_i(x_i) \rightarrow R(x)$, where $R_i \in \text{Rel}_{n_i}(\tau)$, $R \in \text{Rel}_n(\tau)$, $x_i \in V^{n_i}$ and $x \in V^n$, by universal quantification of all their variables.

Similarly, using τ_U and a suitable kind of sentences one can treat (epi-monosource)-topological categories (in the sense of [4]) or semi-topological categories (in the sense of [8]). In either case we get a completion playing the role of Mac Neilles one in the event of topological categories.

A relational theory T of type τ will be called reflexive if for any relation symbol R of τ $T \models (\forall x)R(x, x, \dots, x, \dots)$ holds where $x \in V$. It is transitive if for any cardinal n and any $R \in \text{Rel}_n(\tau)$ $T \models (\forall x) [(\bigwedge_{i \in n} R(x_{i,1}, x_{i,2}, \dots, x_{i,j}, \dots)) \wedge (\bigwedge_{j \in n} R(x_{1,j}, x_{2,j}, \dots, x_{i,j}, \dots))] \rightarrow R(x_{1,1}, x_{2,2}, \dots, x_{i,1}, \dots)]$ holds where $x = (x_{i,j}) \in V^{n \times n}$. Motivating is the case of a binary relation symbol R .

Proposition: Let T be a relational, normal, reflexive and transitive universal Horn theory without equality. Then (\mathcal{A}_T, U_T) is a cartesian closed topological category.

The author conjectures that this proposition can be converted. Namely, one is tempted to seek for a type $\sim \in \tau_U$ such that $(\mathcal{A}_{\sim}, U_{\sim})$ is (in general non-legitimate) cartesian closed topological hull of (\mathcal{A}, U) and its legitimacy corresponds to strict fibre-smallness of (\mathcal{A}, U) in the sense of Adámek and Koubek [2] (i.e. model theoretically

recover their theorem).

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